

Crooked Maps in Finite Fields

Gohar Kyureghyan

*Institute for Algebra and Geometry
Otto-von-Guericke-University Magdeburg
D-39016 Magdeburg*

We consider the maps $f : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$ with the property that the set $\{f(x+a) + f(x) : x \in \mathbb{F}_{2^n}\}$ is a hyperplane or a complement of hyperplane for every $a \in \mathbb{F}_{2^n}^*$. The main goal of the talk is to show that almost all maps $f(x) = \sum_{b \in B} c_b(x+b)^d$, where $B \subset \mathbb{F}_{2^n}$ and $\sum_{b \in B} c_b \neq 0$, are not of that type. In particular, the only such power maps have exponents $2^i + 2^j$ with $\gcd(n, i-j) = 1$. We give also a geometrical characterization of this maps.

Keywords: almost perfect maps, Gold power function, quadrics

1 Introduction

For applications in cryptography the maps $f : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$, which are far from being linear, are important. There are several possibilities to define “being far from linear”. Let $L : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$ be a linear map, then the set $\{L(x+a) + L(x) : x \in \mathbb{F}_{2^n}\}$ consists of only one element $L(a)$ for all fixed $a \in \mathbb{F}_{2^n}$. Hence, a function $f : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$ would be far from being linear if the sets $\{f(x+a) + f(x) : x \in \mathbb{F}_{2^n}\}$ are as big as possible. We can also think about $\{L(x+a) + L(x) : x \in \mathbb{F}_{2^n}\}$ as an affine subspace with only one element, and require for $\{f(x+a) + f(x) : x \in \mathbb{F}_{2^n}\}$ to be one of the largest possible affine subspaces. Another possibility is to require that all coordinate functions $\text{tr}(\alpha f(x)) : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_2$ have large Hamming distance from the linear functions $\text{tr}(\beta x)$ (which are the only linear functions from \mathbb{F}_{2^n} into \mathbb{F}_2). There are three classes of maps with good nonlinearity properties ([4], [1]):

Definition 1 A map $f : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$ is called **almost perfect nonlinear**, if for every $a \in \mathbb{F}_{2^n}^*$

$$|\{f(x+a) + f(x) : x \in \mathbb{F}_{2^n}\}| = 2^{n-1};$$

crooked, if for every $a \in \mathbb{F}_{2^n}^*$

$$\{f(x+a) + f(x) : x \in \mathbb{F}_{2^n}\}$$

is a hyperplane or a complement of a hyperplane;

almost bent, if n is odd and for all $\alpha \in \mathbb{F}_{2^n}^*, \beta \in \mathbb{F}_{2^n}$

$$\mathcal{F}_f(\alpha, \beta) := \sum_{x \in \mathbb{F}_{2^n}} (-1)^{\text{tr}(\alpha f(x) + \beta x)} \in \{-2^{\frac{n+1}{2}}, 0, 2^{\frac{n+1}{2}}\}.$$

Observe, that we extend the notion of crooked maps introduced in [1]. In [1] a map is called crooked if all sets $\{f(x+a) + f(x) : x \in \mathbb{F}_{2^n}\}$ are complements of hyperplanes. These are in our notion bijective crooked maps. Bijective crooked maps exist only for n odd, while crooked maps exist also for n even ([15], [13]). It can be shown ([4], [15], [13]), that

$$\text{crooked} \Rightarrow \text{almost bent} \Rightarrow \text{almost perfect nonlinear.}$$

All known almost perfect nonlinear functions can be obtained from almost perfect nonlinear power maps using the following construction.

Proposition 1 ([3]) *Let $f : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$ be an almost perfect nonlinear map and $l_1, l_2 : \mathbb{F}_{2^n}^2 \rightarrow \mathbb{F}_{2^n}$ be linear maps. Assume that (l_1, l_2) is a permutation on $\mathbb{F}_{2^n}^2$ and $f_2 = l_2(f(x), x)$ is a permutation on \mathbb{F}_{2^n} . Then, the map $f_1 \circ f_2^{-1}$, where $f_1(x) = l_1(f(x), x)$, is almost perfect nonlinear.*

The known exponents of power almost nonlinear maps (up to factor 2^i) are

- $2^k + 1$, $\gcd(k, n) = 1$ (Gold's exponent [9],[1]);
- $2^{2k} - 2^k + 1$, $\gcd(k, n) = 1$ (Kasami's exponent [12])
- $2^{4k} + 2^{3k} + 2^{2k} + 2^k - 1$, if $n = 5k$ (Dobbertin's function [7])
- if $n = 2m + 1$ also
 - $2^m + 3$ (Welch's exponent [6], [2], [11])
 - $2^m + 2^{\frac{m}{2}} - 1$, if m is even, and
 - $2^m + 2^{\frac{3m+1}{2}} - 1$, if m is odd (Niho's exponent [5], [11]);
 - $2^n - 2$ (field inverse [14]).

This list is conjectured to be complete.

The main goal of our talk is to show that the only crooked power maps are the ones with Gold exponents. Denote by C_k the cyclotomic coset modulo $2^n - 1$ containing k , more precisely,

$$C_k = \{k, 2k, \dots, 2^{n-1}k\} \pmod{2^n - 1}.$$

If $|C_k| = l$, then $\{x^k : x \in \mathbb{F}_{2^n}\} \subset \mathbb{F}_{2^l}$ and l is the smallest such number. The binary weight of k is the number of ones in its binary representation. For two integers i and j we write $i \prec j$ if $i \neq j$ and in the binary representations of these integers every digit of i is less or equal to the corresponding digit of j .

We call the integers in the cyclotomic class of $\sum_{j=0}^{\frac{n}{g}-2} 2^{jg}$ exceptional, where g is a divisor of n .

The following results imply that the only crooked power maps are the Gold power maps.

Lemma 1 *Let an integer $0 \leq d \leq 2^n - 2$ have binary weight > 2 and $|C_d| = n$. If for every i with $2^i \prec d$ there exist $j(i)$ and $0 < s(i) < n$ such that $2^{j(i)} \prec d$ and $(d - 2^i) \equiv 2^{s(i)}(d - 2^{j(i)}) \pmod{2^n - 1}$, then d is exceptional.*

Corollary 1 *Let $1 \leq d \leq 2^n - 2$ be an unexceptional integer of binary weight > 2 , $|C_d| = n$. If $f : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$ is given by $f(x) = \sum_{b \in B} c_b(x+b)^d$, where $B \subset \mathbb{F}_{2^n}$ and $\sum_{b \in B} c_b \neq 0$, then the set $\{f(x) + f(x+a) : x \in \mathbb{F}_{2^n}\}$ contains n linearly independent vectors for every $a \in \mathbb{F}_{2^n}^*$.*

Theorem 1 *If $1 \leq d \leq 2^n - 2$ is an unexceptional integer of binary weight > 2 and $|C_d| = n$, then $f(x) = \sum_{b \in B} c_b(x+b)^d$, where $B \subset \mathbb{F}_{2^n}$ and $\sum_{b \in B} c_b \neq 0$, is not crooked.*

In the case $B = \{0\}$ the exceptional exponents can be excluded as well.

Theorem 2 *The only crooked power maps in \mathbb{F}_{2^n} are the ones with exponent $2^i + 2^j$, $\gcd(i - j, n) = 1$.*

It is conjectured [13], that all crooked maps contain only monomials with exponents of binary weight 2 in their polynomial representation. The following observation strengthens this conjecture.

Let n be odd. The almost bent permutations $f(x)$ can be characterized as maps with coordinate functions $\text{tr}(\alpha f(x))$ having the same distances from the hyperplanes as nondegenerate quadrics ([8]). More precisely, let $\alpha, \beta \in \mathbb{F}_{2^n}^*$

$$F_\alpha := \{x \in \mathbb{F}_{2^n} : \text{tr}(\alpha f(x)) = 1\} \text{ and } H_i(\beta) := \{x \in \mathbb{F}_{2^n} : \text{tr}(\beta x) = i\}, \quad i = 0, 1.$$

Then a permutation $f(x)$ is almost bent if and only if

$$F_\alpha \cap H_i(\beta) \in \{2^{n-2}, 2^{n-2} \pm 2^{\frac{n-3}{2}}\}, \quad i = 0, 1,$$

for all $\alpha, \beta \in \mathbb{F}_{2^n}^*$. The following Theorem shows that the coordinate functions of crooked maps behave like quadrics also with the affine subspaces of dimension $n - 2$.

Theorem 3 *Let f be an almost bent permutation with $f(0) = 0$. Then f is crooked if and only if*

$$F_\alpha \cap H_i(\beta_1) \cap H_j(\beta_2) \in \{2^{n-3}, 2^{n-3} \pm 2^{\frac{n-3}{2}}\}, \quad i, j \in \{0, 1\},$$

where $\alpha, \beta_1 \neq \beta_2 \in \mathbb{F}_{2^n}^*$.

Last Theorem was proved using the arguments of the proof for a similar result about power maps in [10].

References

- [1] T. Bending, D. Fon-Der-Flaass, Crooked functions, bent functions, and distance regular graphs, *Electron.J.Comb.* **5** (R34), **14** (1998).
- [2] A. Canteaut, P. Charpin, H. Dobbertin, Binary m-sequences with three-valued crosscorrelation: a proof of Welch's conjecture, *IEEE Trans. Inform. Theory* **46** (2000), 4-8.
- [3] C. Carlet, P. Charpin, V. Zinoviev, Codes, bent functions and permutations suitable for DES-like cryptosystems, *Desings, Codes and Cryptography* **15** (1998), 125-156.
- [4] F. Chabaud and S. Vaudenay, Links between differential and linear cryptanalysis, Advances in Cryptology, EUROCRYPT'94, *Lecture Notes in Computer Science* **950** (1995), 356-365.
- [5] H. Dobbertin, Almost perfectly nonlinear power functions on $\text{GF}(2^n)$: the Niho case, *Information and Computation* **151**, (1999), 57-72.
- [6] H. Dobbertin, Almost perfectly nonlinear power functions on $\text{GF}(2^n)$: the Welch case, *IEEE Transactions on Information Theory* **45** (1999), 1271-1275.
- [7] H. Dobbertin, Almost perfect nonlinear functions on $\text{GF}(2^n)$: a new case for n divisible by 5, in *Finite Fields and Applications*, D. Jungnickel and H. Niederreiter (Eds.), Springer, Berlin 2001, 113-121.

- [8] R.A. Games, The geometry of quadrics and correlations of sequences, *IEEE Trans. Inform. Theory* **32(3)** (1986), 423-426.
- [9] R. Gold, Maximal recursive sequences with 3-valued recursive crosscorrelation functions, *IEEE Trans. Inform. Theory* **14** (1968), 154-156.
- [10] D. Hertel and A. Pott, A characterization of a class of maximum nonlinear functions, submitted.
- [11] H.D.L. Hollmann, Q. Xing, A proof of the Welch and Niho conjectures on cross-correlations of binary m-sequences, *Finite Fields Appl.* **7** (2001), 253-286.
- [12] T. Kasami, The weight enumerators for several classes of subcodes of the second order binary Reed-Muller codes, *Inform. Control.* **18** (1971), 369-394.
- [13] G.M. Kyureghyan, Differentially affine maps, WCC2005, Bergen, Norway pp.296-305 .
- [14] K. Nyberg, Differentially uniform mappings for cryptography, Advances in Cryptology, EURO-CRYPT'93, *Lecture Notes in Computer Science* **765** (1994), 55-64.
- [15] E.R. van Dam and D. Fon-Der-Flaass, Codes, graphs, and schemes from nonlinear functions, *European J. Comb.* **24** (2003), 85-98, doi:10.1016/S0195-6698(02)00116-6.